

Probabilistic Reasoning

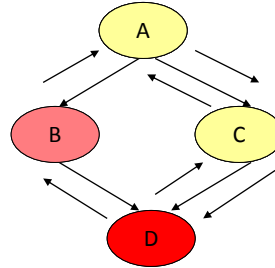
Unit # 5

Intractability of Inference in BN (Laskey)

- **Belief propagation is NP-hard (Cooper, 1988)**
 - NP-hard problems are problems which are at least as computationally complex as NP-complete problems
 - NP-complete is a class of decision problems that are computationally complex
 - There is a theorem which says that if there is an algorithm which solves one NP-Complete problem in polynomial time, then any NP-complete problem can be solved in polynomial time
 - No such algorithm has yet been found, and most people believe there is none
- **Although the general problem appears to be very hard, tractable algorithms are available for certain types of networks**
 - Singly connected networks
 - Some sparsely connected networks
 - Networks amenable to various approximation algorithms

Inference In Multiply Connected Networks

- We get evidence on node D.
- As a result we need to update the marginal probabilities of the rest of the variables
- How to do it?



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3

Exact Inference: Clustering

$P(C)=0.5$

$P(C)=0.5$

C	P(S)
True	0.10
False	0.50

C	P(R)
True	0.8
False	0.20

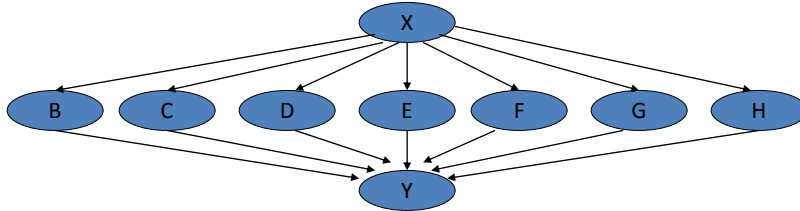
S	R	P(W)
True	True	0.9
True	False	0.8
False	True	0.4
False	False	0.1

S+R				
C	True True	True False	False True	False False
True	0.08	0.02	0.72	0.18
False	0.10	0.40	0.10	0.40

S+R		P(W)
True True	True	0.99
True False	False	0.90
False True	True	0.90
False False	False	0.00

4

Problem with the Clustering Approach



- We are back to the same number of computations as was in the enumeration approach

Overview of Junction/Clique Tree Algorithm (Laskey)

- The basic idea is to transform the graph into clusters of nodes so that the graph of clusters is singly connected and has the *junction tree property* (this property ensures that evidence propagates correctly)
- The junction tree becomes a permanent part of the knowledge representation, and changes only if the graph changes
- Constructing a junction tree from a Bayesian network is inherently non-modular, but the junction tree itself is a modular representation
- Beliefs are propagated in the junction tree using a local message-passing algorithm

How to Construct a JT

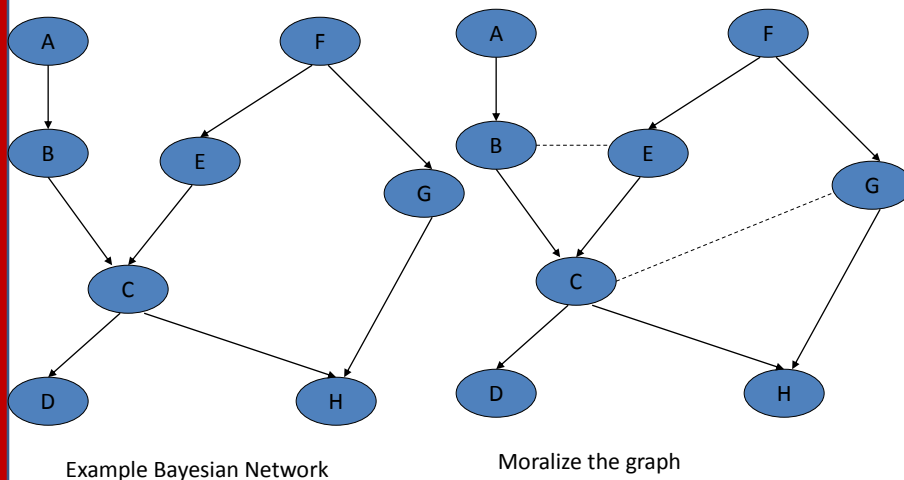
- Moralize the graph by adding links between parents of common children.
- Convert the given BN into an undirected graph by removing the arrowheads.
- Triangulate the graph.
- Order the nodes by maximum cardinality search.
- Find the cliques of the triangulated graph.
- Arrange as a junction tree.

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7

Example I (Laskey)

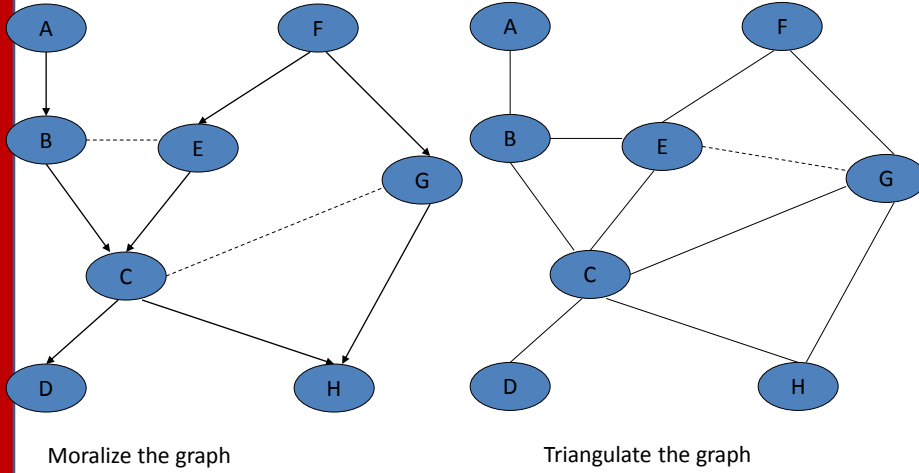


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8

Example I (Cont'd)

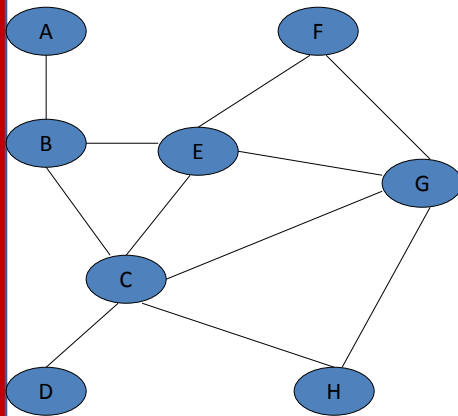


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9

Example I (Cont'd)



Clique	Max. Vertex
AB	2
BEC	4
ECG	5
EGF	6
CGH	7
CD	8

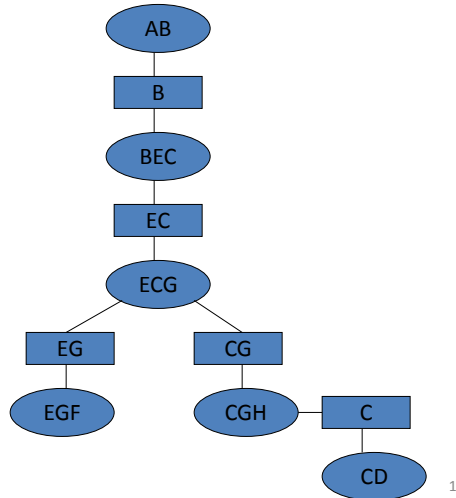
Order Nodes in the Graph

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10

Example I(Cont'd)

Clique	Max. Vertex
AB	2
BEC	4
ECG	5
EGF	6
CGH	7
CD	8



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11

Properties of Clique Tree

- A clique tree is an undirected tree, where each node represents a set of variables, which is called a clique.
- A clique separator is the intersection between its two neighbouring cliques.
- Note that not every clique tree is a junction tree.
- A clique tree is a junction tree if and only if it has the *running intersection property* or the *junction tree property*.

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12

Properties of a Junction Tree (Zhang)

- The junction tree has the following characteristics:
 - it is an undirected tree, its nodes are clusters of variables
 - given two clusters, C_1 and C_2 , every node on the path between them contains their intersection $C_1 \cap C_2$
 - a Separator, S , is associated with each edge and contains the variables in the intersection between neighbouring nodes



Junction Tree Property (Williams)

- A clique tree is a junction tree if it has the following junction tree property:
 - if a node appears in two cliques, it appears everywhere on the path between the cliques.
- *Mathematically, a clique tree possesses the junction tree property if for every pair of cliques V and W , all cliques on the (unique) path between V and W contain $V \cap W$.*
- For every triangulated graph there exists a clique tree which obeys the junction tree property
- Not all clique trees are junction trees

How to Construct a JT

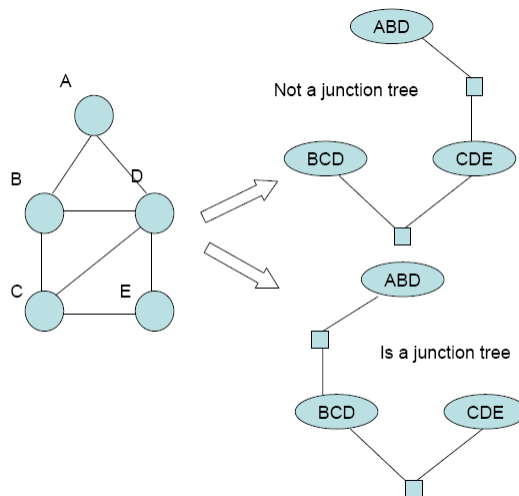
- Moralize the graph by adding links between parents of common children.
- Convert the given BN into an undirected graph by removing the arrowheads.
- Triangulate the graph.
 - If there is a cycle of length > 3 with no chord, add a chord between two non-adjacent vertices in the cycle.
- Order the nodes by maximum cardinality search.
 - Choose any node in the graph and label it 1.
 - For $i=2$ to n
 - Choose the node with the most labeled neighbors and label it i .
 - If any two labeled neighbors of i are not adjacent to each other
 - Add a link between those neighbors
 - Restart the algorithm.
- Find the cliques of the triangulated graph.
- Arrange as a junction tree.

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15

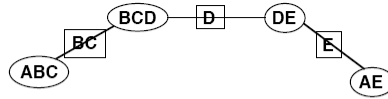
Junction Tree Examples (Jordan)



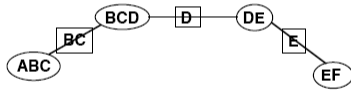
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16

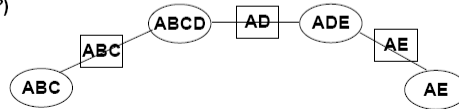
Junction Tree Examples (Laskey)



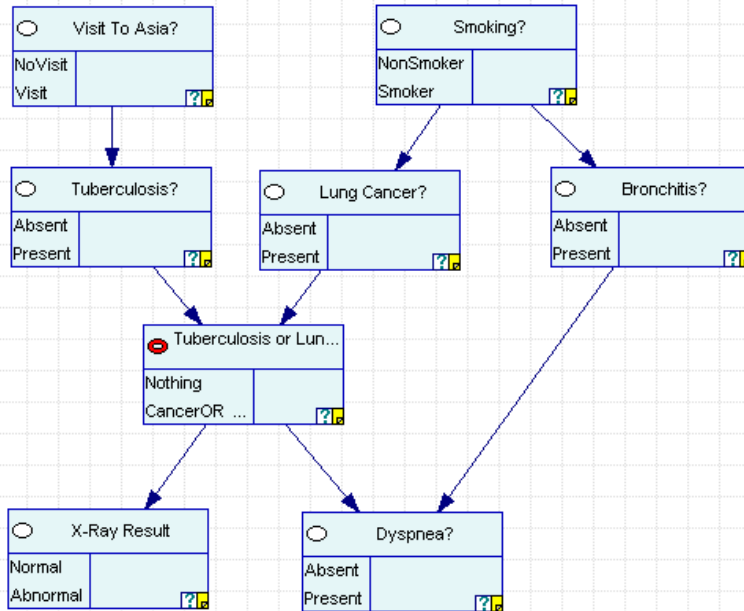
A cluster tree that is not a junction tree
(Why not?)



A junction tree (Why?)

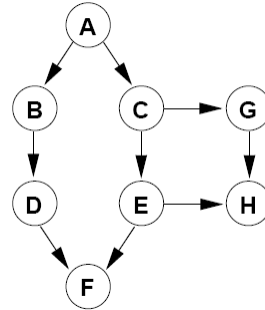


Also a junction tree (Why?)



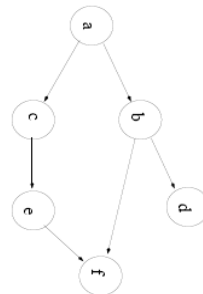
Example III (Huang & Darwiche)

- Consider this BN and transform it into a junction tree by performing moralization, triangulation and maximum cardinality search.

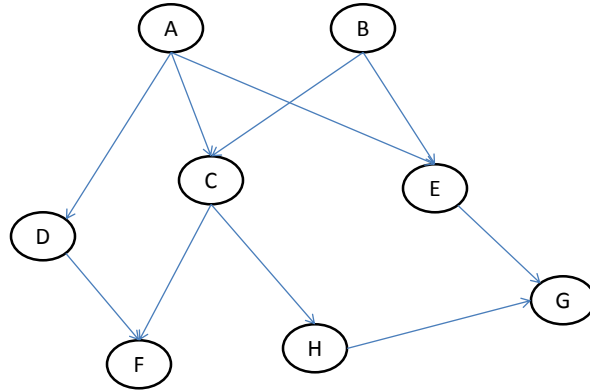


Example IV

- Consider this BN and transform it into a junction tree by performing moralization, triangulation and maximum cardinality search.



Example V (Source: Laskey)

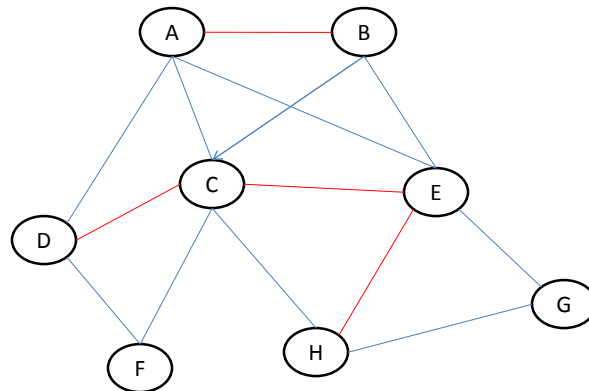


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21

Example V (Cont'd)

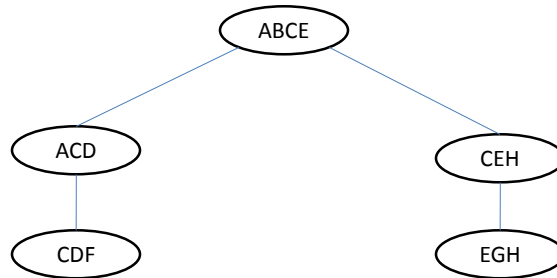


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22

Example V (Cont'd)

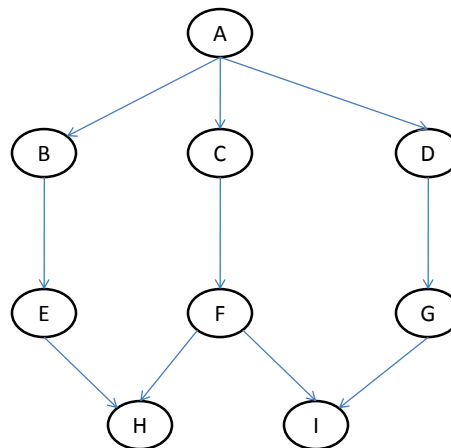


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23

Example VI (Source: Korb & Nicholson)

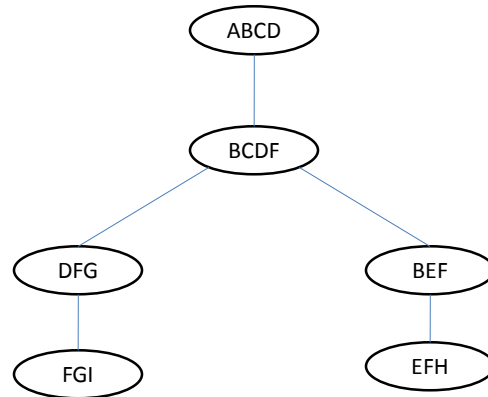


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24

Example VI (Cont'd)



Approximate Inference Mechanisms

- Simulation Based Schemes
 - Logic Sampling
 - Likelihood Weighting
- Approximate Algorithms
 - Loopy Belief Propagation

Logic Sampling (Laskey)

- Logic sampling takes random draws from the joint distribution of all variables in the network
- Each node keeps a running count of how many times each of its states is drawn
- To take one random draw (one observation on all variables):
 - Sample the nodes in order (so parents of a node are sampled before the node)
 - For root node(s), select state with probability $P(\text{state})$
 - For non-root node(s) select state with probability $P(\text{state} | \text{sampled states of parents})$
 - After all nodes have been sampled
 - Check whether sampled values of evidence values match observed states
 - If they don't match, throw away the draw
 - If they match, increment the count for sampled state of each non-evidence node
- Estimate $P(X_t | X_e)$ by observed frequency among the retained samples

Algorithm (From Jensen and Nielsen)

1. Let (X_1, \dots, X_n) be a topological ordering of the variables.
2. For $j = 1$ to N :
 - a) For $i = 1$ to n :
 - Sample a state x_i for X_i using $P(X_i | \text{pa}(X_i) = \pi)$, where π is the configuration already sampled for $\text{pa}(X_i)$.
 - b) If $\mathbf{x} = (x_1, \dots, x_n)$ is consistent with \mathbf{e} , then

$$N(X_k = x_k) := N(X_k = x_k) + 1,$$

where x_k is the state that was sampled for X_k .

3. Return:

$$P(X_k = x_k | \mathbf{e}) \approx \frac{N(X_k = x_k)}{\sum_{x \in \text{sp}(X_k)} N(X_k = x)}.$$

Likelihood Weighting (Laskey)

- Logic sampling can be very inefficient because it throws out so many observations.
- It can be modified to avoid throwing out observations.
- To take one random draw (one observation on all variables):
 - The evidence variables are fixed to their observed states.
 - Generate random samples for the non-evidence variables.
 - Weight the observation by a “sampling weight” that depends upon the evidence.
- Estimate $P(X_t | X_e)$ as the weighted average.

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29

Algorithm (From Jensen and Nielsen)

1. Let (X_1, \dots, X_n) be a topological ordering of the variables.
2. For $j = 1$ to N :
 - a) $w := 1$.
 - b) For $i = 1$ to n :
 - Let \mathbf{x}' be the configuration of (X_1, \dots, X_{i-1}) specified by \mathbf{e} and the previous samples.
 - If $X_i \notin \mathcal{E}$, then:
 - Sample a state x_i for X_i using $P(X_i | \text{pa}(X_i) = \pi)$, where $\text{pa}(X_i) = \pi$ is consistent with \mathbf{x}' .
 - ‡ else
 - $w := w \cdot P(X_i = e_i | \text{pa}(X_i) = \pi)$, where $\text{pa}(X_i) = \pi$ is consistent with \mathbf{x}' .
 - c) $N(X_k = x_k) := N(X_k = x_k) + w$, where x_k is the sampled state for X_k .
3. Return:

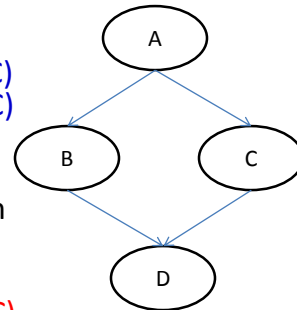
$$P(X_k = x_k | \mathbf{e}) \approx \frac{N(X_k = x_k)}{\sum_{x \in \text{sp}(X_k)} N(X_k = x)}$$

Loopy Belief Propagation

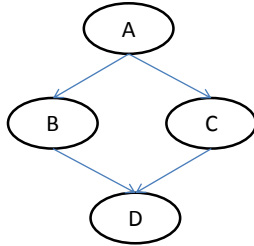
- Originally the message-passing belief propagation was developed for polytrees.
- In loopy belief propagation, we make certain independence assumption (assuming that no loop exists when one exist)

Concept Behind Loopy Belief Propagation

- $P(B) = P(B|A)P(A) + P(B|\sim A)P(\sim A)$
- $P(C) = P(C|A)P(A) + P(C|\sim A)P(\sim A)$
- $P(D) = P(D|B,C)P(B,C) + P(D|B,\sim C)P(B,\sim C) + P(D|\sim B,C)P(\sim B,C) + P(D|\sim B,\sim C)P(\sim B,\sim C)$
- If we assume that B and C are independent (when they are not as obvious from the structure) then we can compute the above equation as
- $P(D) = P(D|B,C)P(B)P(C) + P(D|B,\sim C)P(B)P(\sim C) + P(D|\sim B,C)P(\sim B)P(C) + P(D|\sim B,\sim C)P(\sim B)P(\sim C)$



Examples



SET I

$P(A) = 0.3$
 $P(B|A)=0.7, P(B|\sim A)=0.05$
 $P(C|A)=0.2, P(C|\sim A)=0.9$
 $P(D|B,C) = 0.9, P(D|B,\sim C) = 0.6$
 $P(D|\sim B,C)=0.8, P(D|\sim B,\sim C)=0.02$

SET II

$P(A) = 0.3$
 $P(B|A)=0.9, P(B|\sim A)=0.01$
 $P(C|A)=0.02, P(C|\sim A)=0.95$
 $P(D|B,C) = 0.99, P(D|B,\sim C) = 0.95$
 $P(D|\sim B,C)=0.90, P(D|\sim B,\sim C)=0.02$

- Using Set I and II, compute the probability of D, $P(D)$, using loopy belief propagation and compare the results with the ones obtained through GeNIe (exact, logic sampling and likelihood weighting). What do you observe?